



A generalization of L-fuzzy sets and its application in general system theory

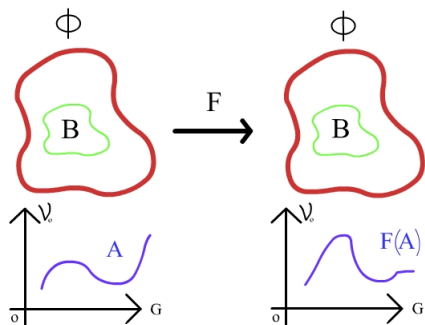
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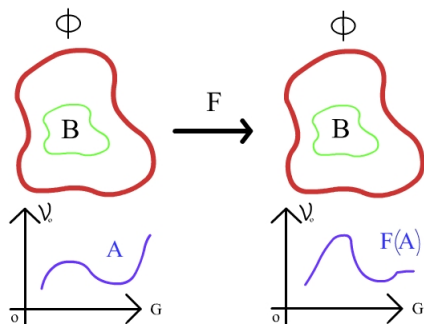


BASIC IDEA





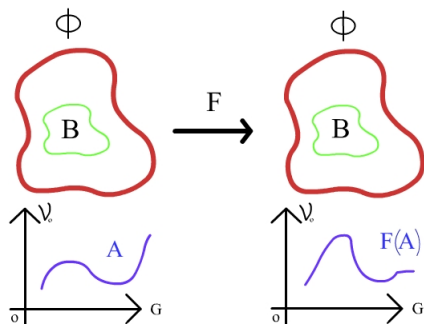
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- ▶ Φ is a **function space** with some nice algebraic and topologic properties (usually **completeness conditions**).



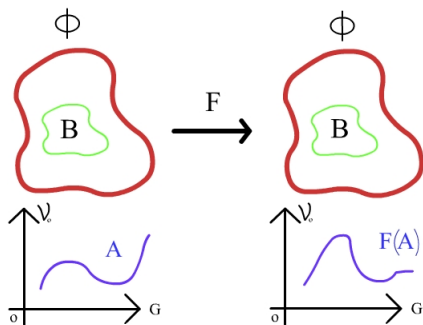
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- Specially Φ is **reconstructable**, i.e. can be reconstructed properly by a (relatively small) subspace B of its **generic** objects.



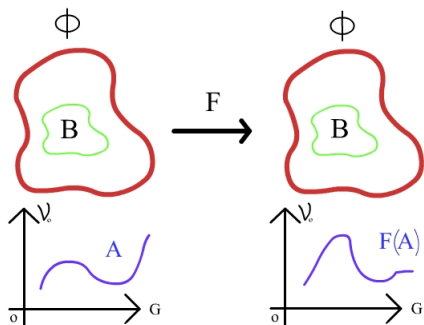
BASIC IDEA



- Usually Φ is chosen to be a product space $G \times \mathcal{V}_0$.
e.g. SPEECH: $G = R$ is the time and $\mathcal{V}_0 = R$ is the space of levels.
e.g. IMAGE: $G = Z^2$ is the two dimensional space $\mathcal{V}_0 = Z$ is the space of gray-scales.



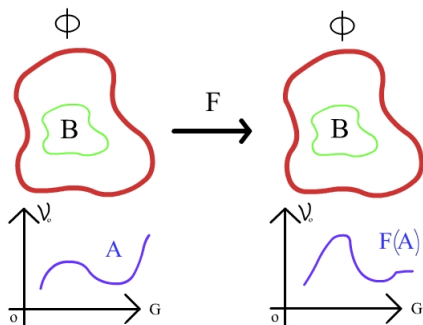
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- F is a **natural map** (i.e compatible with the structure), with nice **representation** properties.



BASIC IDEA



- The whole **setup** should be in coherence with natural phenomena and be able to simulate **input-output** behaviour of such systems (this is called **I/O system theory**).



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- ▶ Φ is a Hilbert space (or some nice Sobolev space).



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$$F(A + \alpha B) = F(A) + \alpha F(B),$$

$$F(T_g(A)) = T_g(F(A)),$$

where T_g is the **translation-by- g operator**, i.e.

$$T_g(A)(t) = A(t - g).$$



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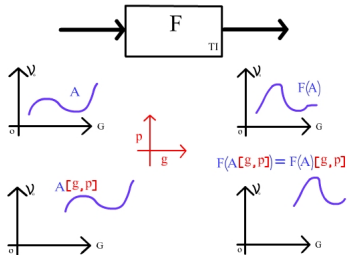
$$T_g(A)(t) = A(t - g).$$

- ▶ We have the following **reconstruction**:

$$F(A)(t) = Conv(A, \delta_F) = \int_{-\infty}^{+\infty} \delta_F(\tau) A(t - \tau) d\tau.$$



TRANSLATION INVARIANCE

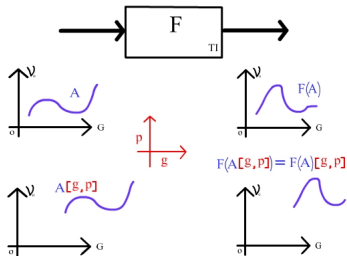


► Consider

$$A = \{(t, A(t)) \mid t \in G\}.$$



TRANSLATION INVARIANCE



- ▶ Then we define the **translation operator** as follows,

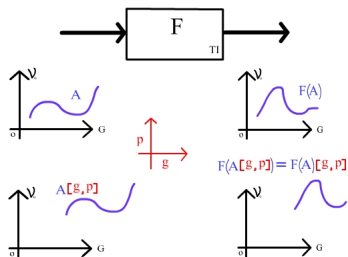
$$A[g, p] = \{(t + g, A(t) * p) \mid t \in G\},$$

which means,

$$A[g, p](t) = A(t - g) * p.$$



TRANSLATION INVARIANCE



- ▶ An operator F is **translation invariant** if

$$F(A[g, p]) = F(A)[g, p].$$

Note: For an LSI operator, translation invariance on the range is equivalent to DC-gain 1.



OUTLINE

Functions v.s. fuzzy sets

Basic morphological operators

The reconstruction property

A deeper approach

Epilogue



OPERATORS V.S. COMPARISON

- ▶ The set $A = \{(t, A(t)) \mid t \in G\} \subseteq G \times \mathcal{V}_0$ can be considered as a **function** or a **fuzzy set**.



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OPERATORS V.S. COMPARISON

- ▶ The set $A = \{(t, A(t)) \mid t \in G\} \subseteq G \times \mathcal{V}_0$ can be considered as a **function** or a **fuzzy set**.
- ▶ The **difference** between the two points of view is in the way we look at the **range** \mathcal{V}_0 .
- ▶ The **functional approach** is when we consider **algebraic properties** and the **fuzzy set approach** is when we consider the **comparative structure** (order structure) of \mathcal{V}_0 .
e.g. The neutral element for summation is 0, while the neutral element for the supremum is $-\infty$.



MINKOWSKI ADDITION AND SUBTRACTION

- Let $(G, +, -, 0)$ be a **group** and $\mathcal{V}_0 = (\Omega, \leq, *, \div, 0) \cup \{-\infty, +\infty\}$ be a **lattice ordered group** with the universal bounds $-\infty$ and $+\infty$ such that,

$$\div(-\infty) = +\infty, \quad \div(+\infty) = -\infty,$$

$$(-\infty) * (+\infty) = (+\infty) * (-\infty) = (+\infty) * (+\infty) = (+\infty)$$

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$$\forall p \in \Omega \quad (+\infty) * p = p * (+\infty) = +\infty,$$

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MINKOWSKI ADDITION AND SUBTRACTION

- Minkowski addition and subtraction for L-fuzzy sets are defined as follows

$$A \oplus B = \sup_g A[g, B(g)] \quad , \quad A \ominus B = \inf_g A[g, \div B(g)].$$



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- ▶ e.g. Let

$$G = \langle a \mid 3a = 0 \rangle,$$

$$A = \{(0, x_1), (a, x_2), (2a, x_3)\},$$

$$B_1 = \{(0, 0), (a, 0), (2a, -\infty)\}.$$

Then,

$$A \oplus B_1 = \{(0, \sup(x_1, x_3)), (a, \sup(x_2, x_1)), (2a, \sup(x_3, x_2))\}.$$



MINKOWSKI EROSION AS CONVOLUTION

- **Minkowski erosion** is defined as follows

$$Er(A, B) = A \ominus B^s.$$

where,

$$B^s = \{(-x, B(x)) \mid x \in G\}.$$



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- ▶ Minkowski erosion is **non-linear** and behaves as a **convolution operator**.



THE KERNEL

- ▶ For a TI operator F the **kernel** is defined as follows

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- ▶ e.g. Let

$$G = \langle a \mid 3a = 0 \rangle,$$

$$A = \{(0, x_1), (a, x_2), (2a, x_3)\},$$

$$F(A) = \{(0, m), (a, m), (2a, m)\}, \text{ where}$$

$$m = \text{Median}(x_1, x_2, x_3).$$

Then,

$$B(F) = \{ \{(0, 0), (a, 0), (2a, -\infty)\}, \{(0, -\infty), (a, 0), (2a, 0)\}, \\ \{(0, 0), (a, -\infty), (2a, 0)\} \}$$



THE RECONSTRUCTION THEOREM

- ▶ A TI operator F is **isotone** if

$$\forall A, B \quad A \leq B \Rightarrow F(A) \leq F(B).$$



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- ▶ **Strong Reconstruction Theorem** (Daneshgar 1995)
Let F be an **isotone TI operator**. Then

$$F(A) = \sup_{D \in K(F)} Er(A, D);$$

and if the base of F exists then

$$F(A) = \sup_{B \in B(F)} Er(A, B).$$



EXAMPLE: THE MEDIAN FILTER

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- ▶ Strong reconstruction theorem implies that
 $\text{Median}(x_1, x_2, x_3) = F(A)(0)$
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- ▶ This is a **fuzzification** of the Boolean expression
 $x_1x_2 + x_2x_3 + x_3x_1$.



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SUM UP

- ▶ **Reconstructions** are usually as **limits of convolutions**.
- ▶ In what **follows** we show that
 - ▶ Such **limits** can be defined in a **variety** of ways.
 - ▶ **Convolutions** are essentially **generalized HOM-functors**.
 - ▶ **Our general setup** will cover both **fuzzy** and **functional** approach.



CATEGORICAL PREREQUISITES I

- ▶ Let $\mathcal{V} = (\mathcal{V}, \cdot, \div, I, a, l, r, c)$ be a **symmetric closed monoidal category** enriched over itself, with identity I , for which $V \cdot - \dashv - \div V$ for any $A \in V$ and the base category \mathcal{V}_0 is complete and cocomplete.



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- ▶ On the other hand, let \mathcal{D} be a class of diagram-schemes and assume that $\mathcal{V}_f \subseteq \mathcal{V}_0$ is a **small full subcategory** such that \mathcal{V}_0 is a **free \mathcal{D} -cocompletion** of \mathcal{V}_f in the sense that,
 - ▶ The totality of all \mathcal{D} -colimits constitute a density presentation for the inclusion $i : \mathcal{V}_f \hookrightarrow \mathcal{V}_0$.
 - ▶ For any $A \in \mathcal{V}_f$, the hom-functor $[A, -]$ preserves all \mathcal{D} -colimits.



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- ▶ Hereafter, $\Delta_W \in \Phi$ is the constant functor with value W .



TWO FUNCTORS

- Consider the following maps for a fixed $D \in \Phi$,

$$\mathbb{T}_D : \text{obj}(\Phi) \longrightarrow \text{obj}(\Phi),$$

$$\mathbb{T}_D(A)_{(z)} = \prod_{|J| < \infty} \left(\bigoplus_{j \in J} W_j \right)^{[A_{(x)}, \bigoplus_j (\Delta_{W_j} \cdot D_{(-x+z)})]_{\Phi}},$$

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where $z \in G$ is a fixed coordinate and \bigoplus is a suitable associative bifunctor.



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where $z \in G$ is a fixed coordinate and \bigoplus is a suitable associative bifunctor.

- **Theorem**(Daneshgar & Hashemi 2000)

Both T_D and H_D can be naturally extended to define **functors** on Φ for any $D \in \Phi$.



AN IMPORTANT SPECIAL CASE

- If \oplus is **preserved** by (\div) (resp. (\cdot)) then

$$\mathsf{T}_D(A)_{(z)} = \prod_{x \in G} A_{(x)} \cdot D_{(-x+z)} \stackrel{\text{def}}{=} \dot{\mathsf{T}}_D(A)_{(z)}$$

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For any $D \in \Phi$ we have $\dot{\mathsf{T}}_D \dashv \dot{\mathsf{H}}_D$ (i.e. $\dot{\mathsf{T}}_D$ is the **left adjoint** of $\dot{\mathsf{H}}_D$).



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- ▶ **Theorem**(Daneshgar& Hashemi 2000)

There exist natural isomorphisms $\tilde{a}, \tilde{l}, \tilde{r}$ and \tilde{c} such that for any $D \in \text{obj}(\Phi)$, $\dot{\Phi} = (\Phi, \dot{\mathsf{T}}_D, \dot{\mathsf{H}}_D, P^{0,I}, \tilde{a}, \tilde{l}, \tilde{r}, \tilde{c})$ is a **symmetric closed monoidal category**.



GENERAL RECONSTRUCTION THEOREM

- ▶ A **point** $P^{g,V} \in \Phi$ with value $V \in \text{obj}(\mathcal{V}_0)$ at coordinate $g \in G$ is defined (up to isomorphism) as

$$P^{g,V}(x) \stackrel{\text{def}}{=} \begin{cases} V & x = g \\ -\infty & x \neq g. \end{cases}$$



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- ▶ **Theorem**(Daneshgar& Hashemi 2000)
Let $F : \dot{\Phi} \longrightarrow \dot{\Phi}$ be a $\dot{\Phi}$ -functor such that F **preserves** the internal Hom of $\dot{\Phi}$, \dot{H}_D , for any $D \in \dot{\Phi}$; and $\mathit{Ker}(F)$ is a **\mathcal{D} -type diagram-scheme**. Then F has a representation as a \mathcal{D} -colimit of representables as

$$F(A) \simeq \text{Colim}_{(D,d) \in \mathit{Ker}(F)} \dot{H}_D(A).$$



EXAMPLE I

- ▶ Consider $G = (\mathbf{R}, +)$ and $\mathcal{V} = (\mathbf{R}, +, \leq)$ with two universal bounds $+\infty$ and $-\infty$. Then \dot{T} is the **Minkowski addition** and \dot{H} is the **erosion operator**.



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- ▶ Also, it can be shown that in this case being a $\dot{\Phi}$ -functor is **equivalent** to the definition of a **translation invariant** operator, where we have

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- ▶ **Note** that in this case the general reconstruction theorem yields the classical reconstruction theorem for TI operators.



THE UNIFORM CASE

- Now consider the following maps for a fixed $D \in \Phi$,

$$\tilde{T}_D : \Phi \longrightarrow \mathcal{V},$$

$$\tilde{T}_D(A) = \prod_{|J| < \infty} \left(\bigoplus_{j \in J} W_j \right)^{[A(x), \bigoplus_j (\Delta_{W_j} \cdot D_{(-x+z_j)})]_{\Phi}},$$

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where \bigoplus is a suitable associative bifunctor.



A SIMPLE CASE

- As a **simple** case we have,

$$\ddot{T}_D(A) = \prod_j W_j^{[A(x), \Delta_{W_j} \div D(-x+z_j)]_\Phi} = \prod_z \dot{T}_D(A)_{(z)},$$

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- ▶ **Theorem**(Daneshgar& Hashemi 2000)

If the tensor of \mathcal{V} is **preserved by coproducts**, then there exist a natural composition law and an identity map such that

$[D, B] \stackrel{\text{def}}{=} \ddot{H}_D(B)$, as the internal Hom, turns Φ into a \mathcal{V} -category.



A SIMPLE CASE

- ▶ As a **simple** case we have,

$$\ddot{T}_D(A) = \prod_j W_j^{[A(x), \Delta_{W_j} \div D(-x+z_j)]\Phi} = \prod_z \dot{T}_D(A)_{(z)},$$

and

$$\ddot{H}_D(B) = \prod_j V_j^{[\Delta_{V_j} \cdot D(-z_j+x), B(x)]\Phi} = \prod_z \dot{H}_D(A)_{(z)}.$$

- ▶ **Theorem**(Daneshgar& Hashemi 2000)

If the tensor of \mathcal{V} is **preserved by coproducts**, then there exist a natural composition law and an identity map such that

$[D, B] \stackrel{\text{def}}{=} \ddot{H}_D(B)$, as the internal Hom, turns Φ into a \mathcal{V} -category.

- ▶ **Can you prove a similar theorem for the general case?**



EXAMPLE II

- Let $G = (\mathbf{R}, +)$ and $\mathcal{V} = (\mathbf{R}^+, \cdot, \geq)$. Then

$$\ddot{H}_D(B) = \inf_z (\sup_x (B_{(x)} \div D_{(-z+x)}))$$

in the ordinary order of real numbers.



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- This is the **Haar fraction!**



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- ▶ In the general setting a **closed monoidal category** with suitable **completeness** properties is a suitable model for the valuation space that generalizes **both** functional and fuzzy approaches.



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- ▶ The **fuzzy set approach** introduces applications based on **Boolean functions** that give rise to **non-linear** operators.
- ▶ Can one define a **limit process** that simulates **integration** in the general setting?
- ▶ Can one develop **transformation techniques** that facilitate a design process based on some **wanted** behaviour of operators. (This should be definitely related to some **new understandings and dimensions in system design.**)



THE END.

Thank You!